Rule-Based Composite Event Queries: The Language XChange^{EQ} and its Semantics

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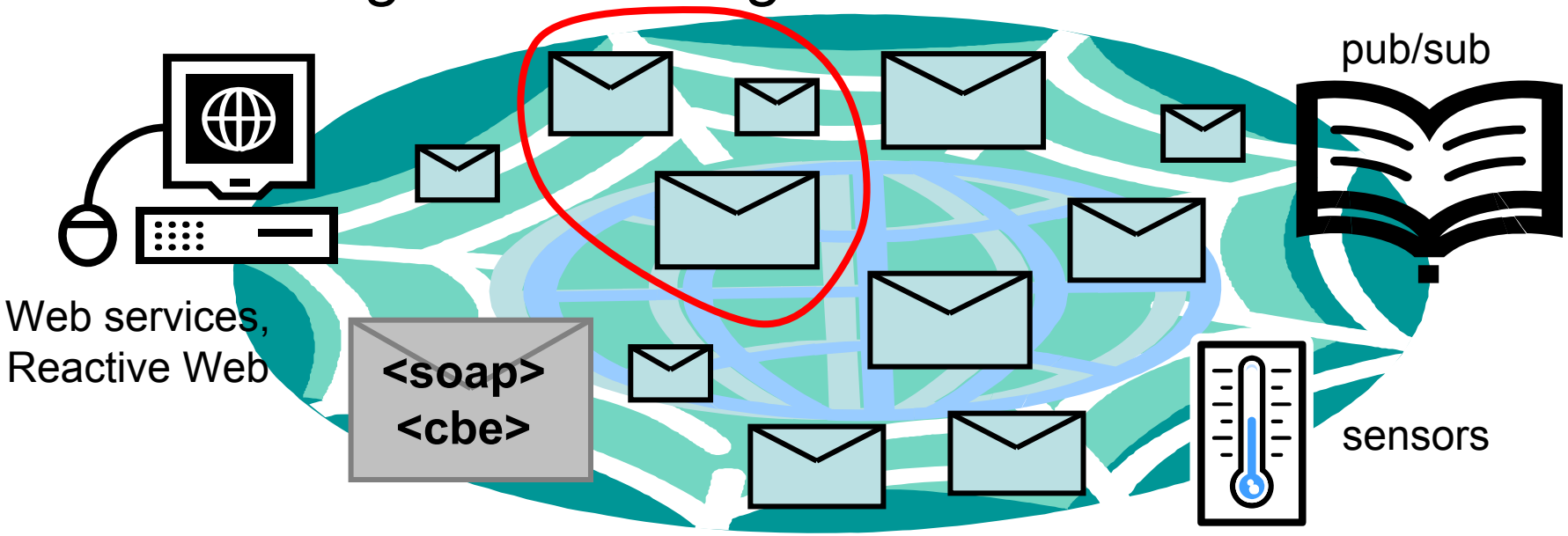
Motivation: Composite Events

Generating and reacting to events on the Web

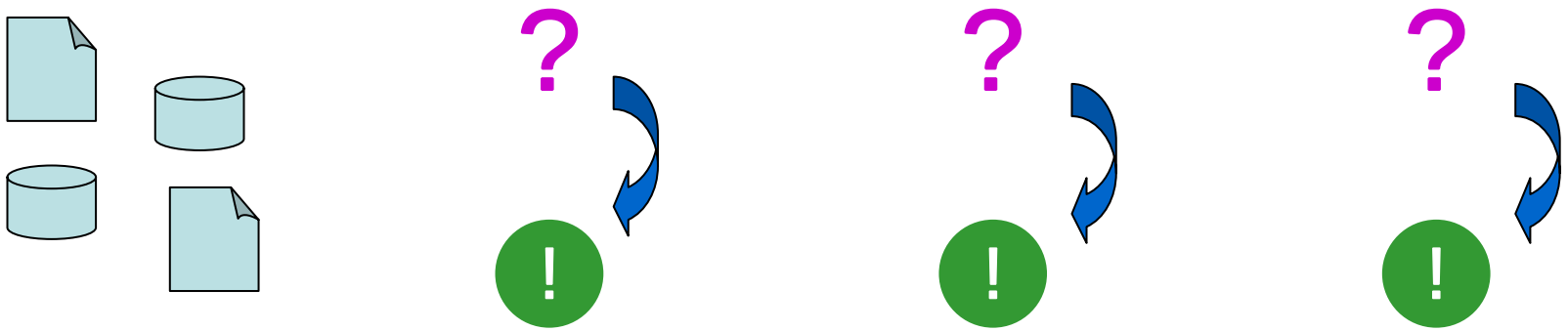


Composite Events

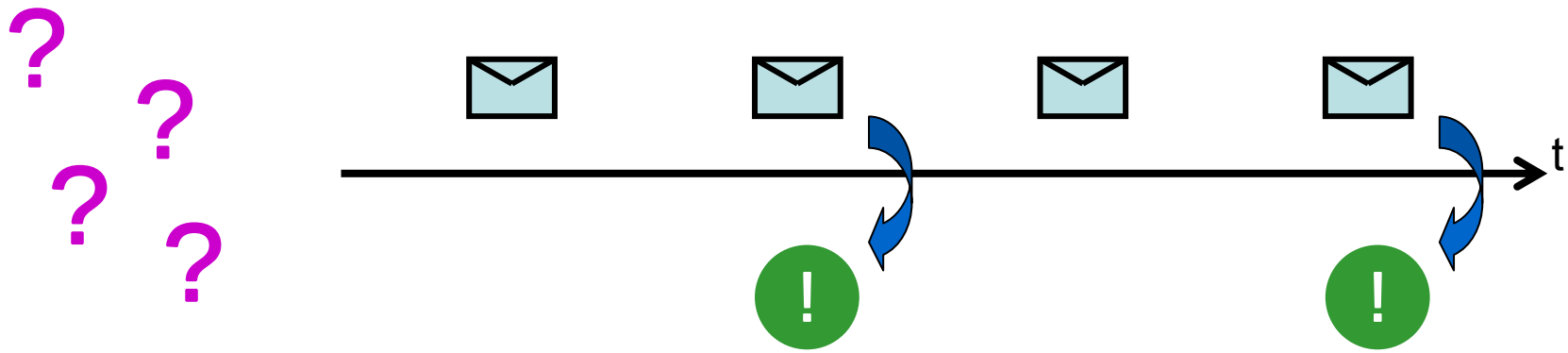
- Must be inferred from “atomic events” (messages)
- Multiple atomic events, relationship between them
- Need query language!

Queries against Event Streams

Database Queries, Web Queries:

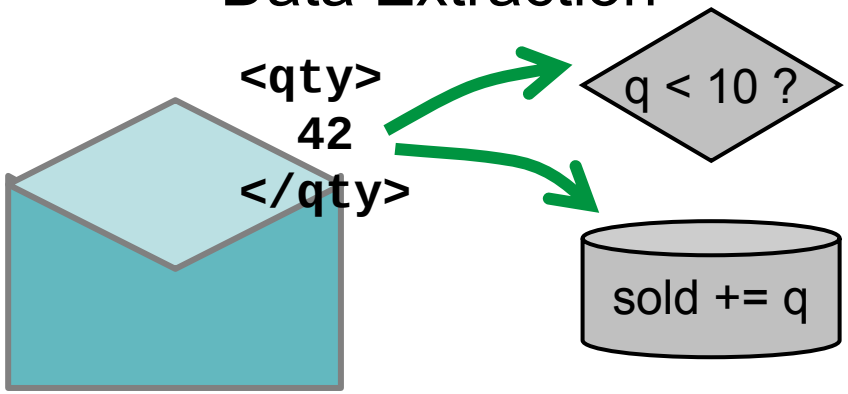


(Composite) Event Queries:

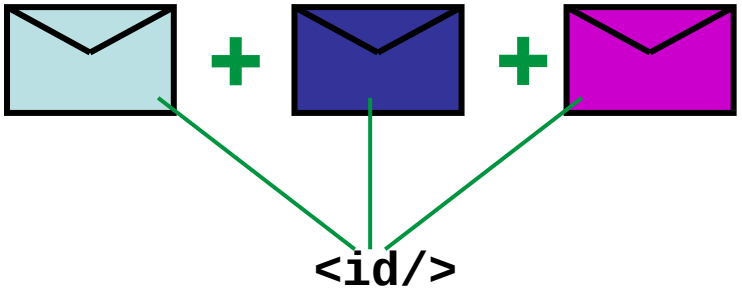


Language Requirements

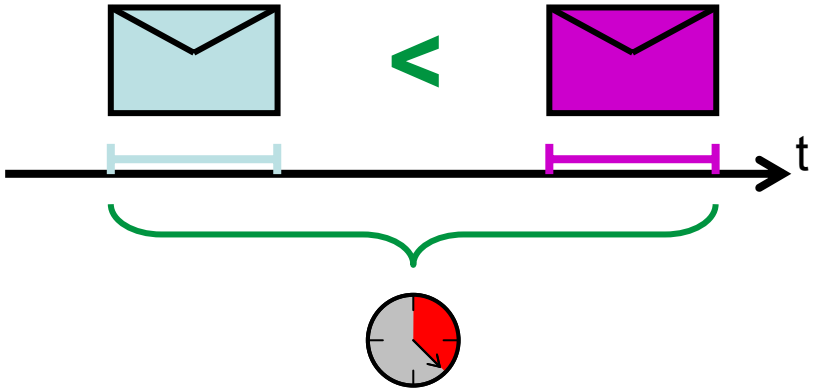
Data Extraction



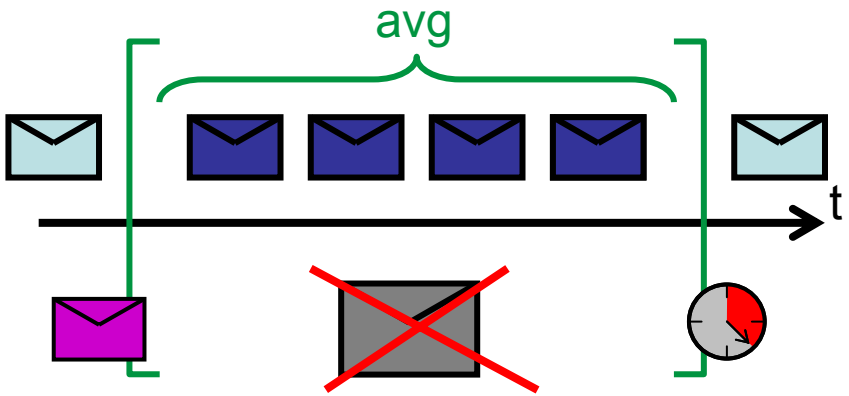
Event Composition



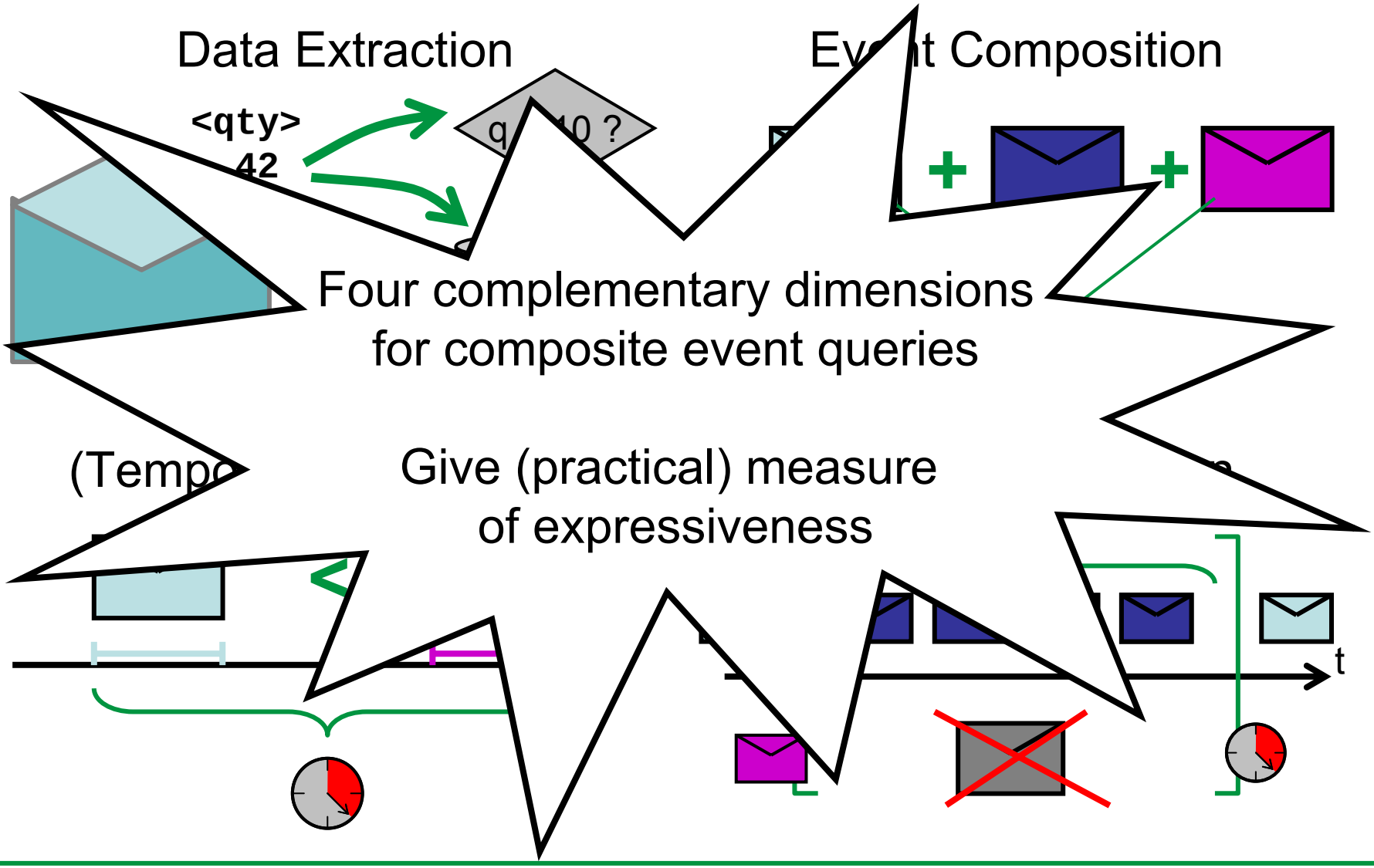
(Temporal) Relationships



Event Accumulation



Language Requirements



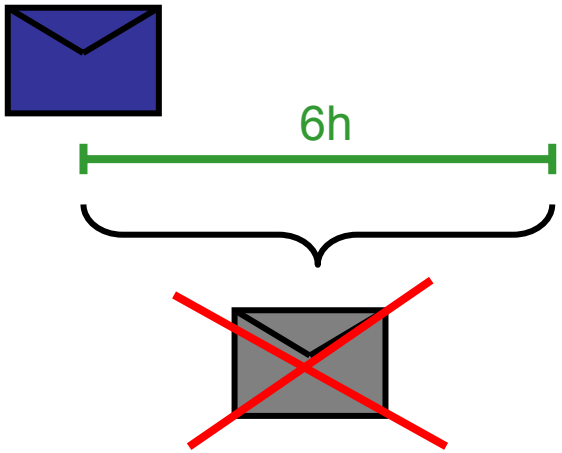
XChange^{EQ}: Rule-Based CEQs

- High-level, declarative query language for composite events, fully covers four dimensions
- Pattern-based queries on XML event messages: embeds Web query language Xcerpt
- Integrates into reactive rule language XChange
  - Perform automatic reactions, timing important
- Deductive (event) rules:
  - Define new, “virtual” events from received events
  - mediation, abstraction, reasoning (cf. database views)
  - Side-effect free; don’t implement by reactive rules: optimization, (human) understanding

XChange^{EQ}: Example Rule

```
DETECT
  overdue {
    id { var ID }
    cust { var C } }
ON
  and {
    event o: order {{
      id { var ID },
      quantity { var Q },
      cust { var C } }}
    event w: extend[o, 6h],
    while w: not shipped {{
      id { var ID } }}
  } where { var Q < 10 }
END
```

```
DETECT
  order {
    id { var ID},
    quantity { var C },
    cust { var C } }
FROM
  ...
END
```



Semantics (1)

- Declarative Semantics for XChangeEQ:
  model + fixpoint theories for stratified programs
  (A standard approach for rule languages)
- (Tarski-style) model theory:
  Define M $\models$ F^t recursively

$$
\begin{array}{ll}
I, \Sigma, \tau \models (\texttt{event } i : q)^t & \text{iff exists } e^{t'} \in I \text{ with } \tau(i) = e^{t'},\ t' = t, \\
 & \text{and for all } e' \in \Sigma(q) \text{ we have } e' \preceq e \\
I, \Sigma, \tau \models (\texttt{event } i : \texttt{extends}[j,d])^t & \text{iff exists } e^{t'} \text{ with } \tau(j) = e^{t'},\ \tau(i) = e^t,\ t = t' + d \\
\ldots & \text{(Definitions for other temporal events are similar and skipped.)} \\
M \models (q_1 \wedge q_2)^t & \text{iff } M \models q_1^{t_1} \text{ and } M \models q_2^{t_2} \text{ and } t = t_1 \sqcup t_2 \\
M \models (q_1 \vee q_2)^t & \text{iff } M \models q_1^t \text{ or } M \models q_2^t \\
I, \Sigma, \tau \models (Q \texttt{ where } C)^t & \text{iff } I, \Sigma, \tau \models Q^t \text{ and } W_{\Sigma,\tau}(C) = true \\
I, \Sigma, \tau \models (\texttt{while } j : \texttt{ not } q)^t & \text{iff exists } e^{t'} \text{ with } \tau(j) = e^{t'},\ t' = t, \\
 & \text{and for all } t'' \sqsubseteq t \text{ we have } I, \Sigma, \tau \not\models q^{t''} \\
I, \Sigma, \tau \models (\texttt{while } j : \texttt{ collect } q)^t & \text{iff exists } e^{t'} \text{ with } \tau(j) = e^{t'},\ t' = t, \text{ and exist } n \geq 0, \\
 & \Sigma_1, \ldots \Sigma_n,\ t_1 \sqsubseteq t, \ldots t_n \sqsubseteq t \text{ with } \Sigma = \bigcup_{i=1..n} \Sigma_i, \\
 & \text{and for all } i = 1..n \text{ we have } I, \Sigma_i, \tau \models q^{t_i} \\
I, \Sigma, \tau \models (c \leftarrow Q)^t & \text{iff (1) } \Sigma'(c)^t \subseteq I \text{ for } \Sigma' \text{ maximal (w.r.t. } FreeVars(Q)) \text{ and } \tau' \\
 & \text{such that } I, \Sigma', \tau' \models Q^t, \text{ or (2) } I, \Sigma', \tau' \not\models Q^t \text{ for all } \Sigma', \tau'
\end{array}
$$

$$
\begin{array}{ll}
W_{\Sigma,\tau}(i \texttt{ before } j) = \textbf{true} & \text{iff } end(\tau(i)) < begin(\tau(j)) \\
W_{\Sigma,\tau}(i \texttt{ during } j) = \textbf{true} & \text{iff } begin(\tau(j)) < begin(\tau(i)) \text{ and } end(\tau(i)) < end(\tau(j)) \\
W_{\Sigma,\tau}(i \texttt{ overlaps } j) = \textbf{true} & \text{iff } begin(\tau(j)) < begin(\tau(i)) < end(\tau(j)) < end(\tau(i))
\end{array}
$$

- Accommodates event identifiers (“event o:”)
- Events have occurrence times
- Temporal relations: fixed interpretation

Semantics (2)

- Restriction to stratified programs
  - w.r.t. negation, grouping, *relative temporal events*
- Fixpoint: model $M_{P,E}$
  - $T_P(I)$: all events derivable by rules in P
  - starting with incoming event stream E
  - compute fixpoints stratum by stratum

$$P = P_1 \uplus \cdots \uplus P_n$$

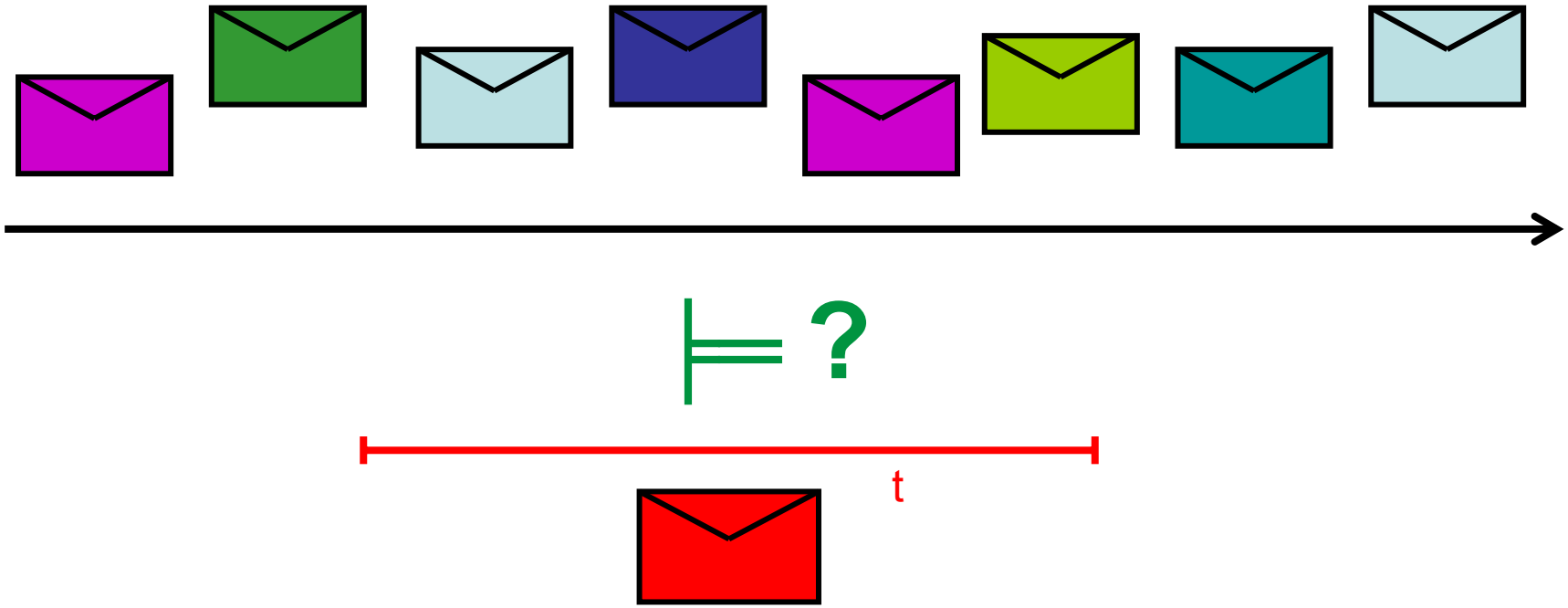
$$T_P(I) = I \cup \{e^t \mid \text{there exist a rule } c \leftarrow Q \in P, \text{ a maximal substitution set } \Sigma, \text{ and a substitution } \tau \text{ s.t. } I, \Sigma, \tau \models Q^t \text{ and } e \in \Sigma(c)\}$$

T_P^ω: least fixpoint of T_P

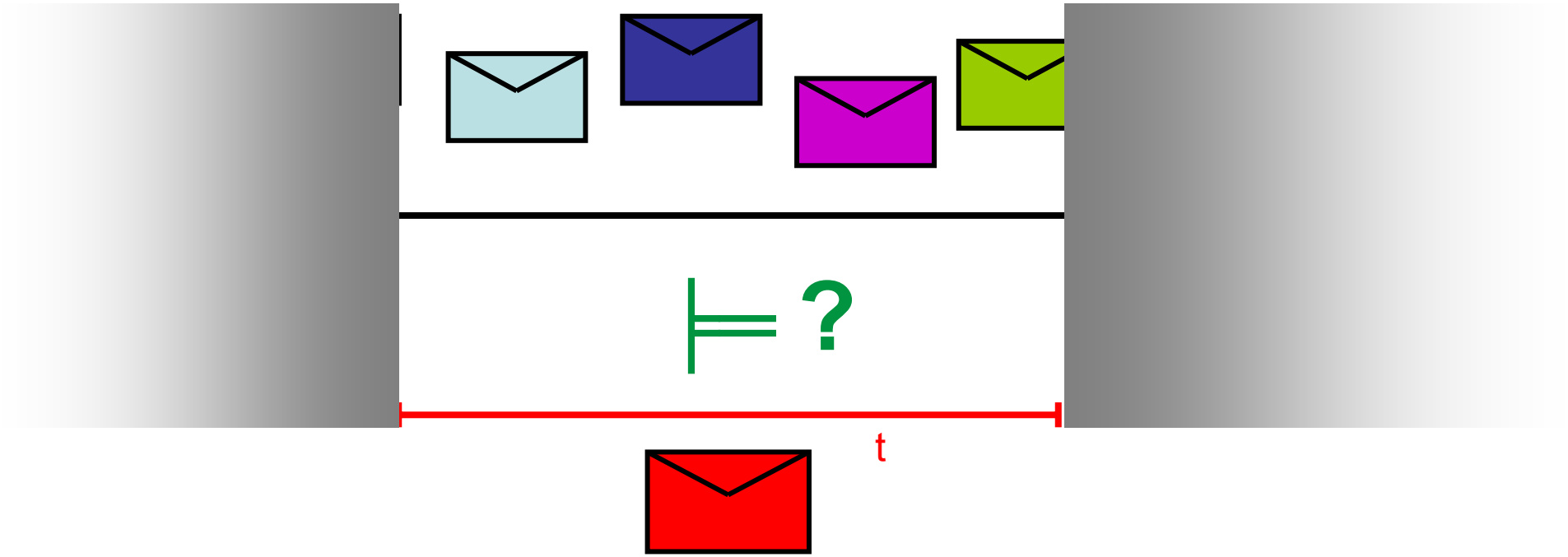
$$M_0 = E = T_\emptyset^\omega(E),$$
$$M_1 = T_{P_1}^\omega(M_0),$$
$$\ldots,$$
$$M_n = T_{P_n}^\omega(M_{n-1}) =: M_{P,E}.$$

- Theorem:
  - P stratified program, E (incoming) event stream.
  - Then: $M_{P,E}$ is a minimal model of P under E and
  - Independent of the stratification of P

Unbounded Event Streams



Unbounded Event Streams



- XChange^{EQ}-Semantics are well-defined for *unbounded* (“infinite”) incoming event streams E

$$M_{P,E} \mid t = M_{P,E|t} \mid t$$

Summary and Outlook

- XChange^{EQ}:
  - High-level event query language
  - Full coverage of all four dimensions, XML support
  - Support for (deductive) event rules
- Declarative Semantics
  - Model and fixpoint theory for stratified programs
  - Well-defined on unbounded event streams
- Outlook
  - Incremental, data-driven evaluation
  - Optimizations based on temporal conditions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